

NOTE

A Short Proof of Mader's \mathcal{S} -Paths Theorem

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Received April 19, 2000

For an undirected graph $G = (V, E)$ and a collection \mathcal{S} of disjoint subsets of V , an \mathcal{S} -path is a path connecting different sets in \mathcal{S} . We give a short proof of Mader's min-max theorem for the maximum number of disjoint \mathcal{S} -paths. © 2001 Academic Press

Let $G = (V, E)$ be an undirected graph and let \mathcal{S} be a collection of disjoint subsets of V . An \mathcal{S} -path is a path connecting two different sets in \mathcal{S} . Mader [4] gave the following min-max relation for the maximum number of (vertex-) disjoint \mathcal{S} -paths, where $S := \bigcup \mathcal{S}$.

MADER'S \mathcal{S} -PATHS THEOREM. *The maximum number of disjoint \mathcal{S} -paths is equal to the minimum value of*

$$|U_0| + \sum_{i=1}^n \lfloor \frac{1}{2} |B_i| \rfloor, \quad (1)$$

taken over all partitions U_0, \dots, U_n of V such that each \mathcal{S} -path disjoint from U_0 traverses some edge spanned by some U_i . Here B_i denotes the set of vertices in U_i that belong to S or have a neighbour in $V \setminus (U_0 \cup U_i)$.

Lovász [3] gave an alternative proof by deriving it from his matroid matching theorem. Here we give a short proof of Mader's theorem.

Let μ be the minimum value obtained in (1). Trivially, the maximum number of disjoint \mathcal{S} -paths is at most μ , since any \mathcal{S} -path disjoint from U_0 and traversing an edge spanned by U_i traverses at least two vertices in B_i .

I. First, the case where $|T| = 1$ for each $T \in \mathcal{S}$ was shown by Gallai [2] by reduction to matching theory as follows: Let the graph $\tilde{G} = (\tilde{V}, \tilde{E})$ arise from G by adding a disjoint copy G' of $G - S$ and making the copy v' of each $v \in V \setminus S$ adjacent to v and to all neighbours of v in G .



We claim that \tilde{G} has a matching of size $\mu + |V \setminus S|$. Indeed, by the Tutte-Berge formula [5, 1] it suffices to prove that for any $\tilde{U}_0 \subseteq \tilde{V}$,

$$|\tilde{U}_0| + \sum_{i=1}^n \lfloor \frac{1}{2} |\tilde{U}_i| \rfloor \geq \mu + |V \setminus S|, \quad (2)$$

where $\tilde{U}_1, \dots, \tilde{U}_n$ are the components of $\tilde{G} - \tilde{U}_0$. Now if for some $v \in V \setminus S$ exactly one of v, v' belongs to \tilde{U}_0 , then we can delete it from \tilde{U}_0 , thereby not increasing the left-hand side of (2). So we can assume that for each $v \in V \setminus S$, either $v, v' \in \tilde{U}_0$ or $v, v' \notin \tilde{U}_0$. Let $U_i := \tilde{U}_i \cap V$ for $i = 0, \dots, n$. Then U_1, \dots, U_n are the components of $G - U_0$, and we have

$$|\tilde{U}_0| + \sum_{i=1}^n \lfloor \frac{1}{2} |\tilde{U}_i| \rfloor = |U_0| + \sum_{i=1}^n \lfloor \frac{1}{2} |U_i \cap S| \rfloor + |V \setminus S| \geq \mu + |V \setminus S| \quad (3)$$

(since in this case $B_i = U_i \cap S$ for $i = 1, \dots, n$), showing (2).

So \tilde{G} has a matching M of size $\mu + |V \setminus S|$. Let N be the matching $\{vv' \mid v \in V \setminus S\}$ in \tilde{G} . As $|M| = \mu + |V \setminus S| = \mu + |N|$, the union $M \cup N$ has at least μ components with more edges in M than in N . Each such component is a path connecting two vertices in S . Then contracting the edges in N yields μ disjoint \mathcal{S} -paths in G .

II. We now consider the general case. Fixing V , choose a counterexample E, \mathcal{S} minimizing

$$|E| - |\{ \{t, u\} \mid t, u \in V, \exists T, U \in \mathcal{S} : t \in T, u \in U, T \neq U \}|. \quad (4)$$

By Part I, there exists a $T \in \mathcal{S}$ with $|T| \geq 2$. Then T is independent in G , since any edge e spanned by T can be deleted without changing the maximum and minimum value in Mader's theorem (as any \mathcal{S} -path traversing e contains an \mathcal{S} -path not traversing e and as deleting e does not change any set B_i), while decreasing (4).

Choose $s \in T$. Replacing \mathcal{S} by $\mathcal{S}' := (\mathcal{S} \setminus \{T\}) \cup \{T \setminus \{s\}, \{s\}\}$ decreases (4), but not the minimum in Mader's theorem (as each \mathcal{S} -path is an \mathcal{S}' -path and as $\bigcup \mathcal{S}' = S$). So there exists a collection \mathcal{P} of μ disjoint \mathcal{S}' -paths. We can assume that no path in \mathcal{P} has any internal vertex in S .

Necessarily, there is a path $P_0 \in \mathcal{P}$ connecting s with another vertex in T , all other paths in \mathcal{P} being \mathcal{S} -paths. Let u be an internal vertex of P_0 . Replacing \mathcal{S}' by $\mathcal{S}'' := (\mathcal{S}' \setminus \{T\}) \cup \{T \cup \{u\}\}$ decreases (4), but not the minimum in Mader's theorem (as each \mathcal{S}' -path is an \mathcal{S}'' -path and as $\bigcup \mathcal{S}'' \supset S$). So there exists a collection \mathcal{Q} of μ disjoint \mathcal{S}'' -paths. Choose \mathcal{Q} such that no internal vertex of any path in \mathcal{Q} belongs to $S \cup \{u\}$ and such that \mathcal{Q} uses a minimal number of edges not used by \mathcal{P} .

Necessarily, u is an end of some path $Q_0 \in \mathcal{Q}$, all other paths in \mathcal{Q} being \mathcal{S} -paths. As $|\mathcal{P}| = |\mathcal{Q}|$ and as u is not an end of any path in \mathcal{P} , there exists

an end v of some path $P \in \mathcal{P}$ that is not an end of any path in \mathcal{Q} . Now P intersects at least one path in \mathcal{Q} (since otherwise $P \neq P_0$, and $(\mathcal{Q} \setminus \{Q_0\}) \cup \{P\}$ would consist of μ disjoint \mathcal{S} -paths). So when following P starting at v , there is a first vertex w that is on some path in \mathcal{Q} , say, on $Q \in \mathcal{Q}$.

For any end x of Q let Q^x be the $x-w$ part of Q . Let P^v be the $v-w$ part of P and let U be the set in \mathcal{S}'' containing v . Then for any end x of Q we have that Q^x is part of P or the other end of Q belongs to U , since otherwise by rerouting part Q^x of Q along P^v , Q remains an \mathcal{S}'' -path disjoint from the other paths in \mathcal{Q} , while we decrease the number of edges used by \mathcal{Q} and not by \mathcal{P} , contradicting the minimality assumption.

Let y, z be the ends of Q . We can assume that $y \notin U$. Then Q^z is part of P , hence Q^y is not a part of P (as Q is not a part of P , as otherwise $Q = P$, and hence v is an end of Q), so $z \in U$. As z is on P and also as v belongs to U and is on P , we have $P = P_0$. So $U = T \cup \{u\}$ and $Q = Q_0$ (since Q^z is part of P , so $z = u$). But then rerouting part Q^z of Q along P^v gives μ disjoint \mathcal{S} -paths, contradicting our assumption.

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